

Network Traffic Processing in Distributed Environment

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Abstract

The growth of computer networks and the Internet availability opens new opportunities for cybercrime activities. Security administrators and LEA (Law Enforcement Agency) officers call for powerful tools for high-speed network communication analysis of an enormous amount of traffic. The forensic analysis needs for various cybercrime cases may differ. This paper aims to design a novel approach of real-time network traffic processing up to an application layer in a distributed environment. The research focuses on captured traffic analysis and information extraction of multiple application protocols. The solution has to be configurable, scalable and capable to analyze even incomplete communication.

Keywords: Network forensic analysis — Network traffic processing — Actor model

Supplementary Material: [Demonstration Video](#)

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1. Introduction

1 The expansion of computer networks and the Internet
2 availability opens new opportunities for cybercrime ac-
3 tivities and security incidents associated with network
4 applications. The amount of connected devices grows,
5 and the traffic speed increases. Security administrators
6 and LEA (Law Enforcement Agency) officers call for
7 powerful tools that enables them to extract useful in-
8 formation from network communication. The network
9 forensics that is responsible for capturing, collecting
10 and network data analyzing is getting more important.
11

12 In the forensic investigation, the network traffic
13 is continuously captured from multiple sources. The
14 captured data has to be processed and analyzed up
15 to the application layer. For LEA, the interesting in-
16 formation from computer network traffic is primarily
17 hidden in application messages such as instant mes-
18 saging, emails, voice, RTP, localizable information,

documents, pictures, etc. Although, the relevance of 19
extracted artifacts may be different from case to case, 20
all cases require at least some kind of digital evidence. 21
The processing system has to be able to extract this 22
data from the traffic, even if it is corrupted. 23

The analysis of high-speed traffic is easier to achieve 24
in a distributed environment. I have decided to use 25
actor model that is effective, and capable of linear 26
scalability. Scalable properties of *actor model* design 27
for network forensics are promising as the VAST plat- 28
form [1] shown. 29

In this work, I intend to design such system as 30
a software solution that is also linearly scalable, and 31
platform independent. In comparison to VAST, I want 32
to perform real-time analysis as well with the focus on 33
information extraction from the application layer. 34

2. Background & Related Work

Network forensics is a process that identifies, captures and analyzes network traffic [2]. Network forensic techniques are used by several network forensic frameworks [3, 4, 5, 6, 7, 8] and tools. In this research apart of framework, I have examined also open source tools. PyFlag and CapAnalysis are capable of offline analysis, as well as Xplico and NetworkMiner that can also perform real-time analysis. All mentioned tools do not support distributed deployment.

The models for distributed processing [9, 10, 11] are more suitable for real-time network forensic analysis from multiple sources, such as logs and captured communication. The models are based on an agent system, where numerous agents performs collection task. The extracted information is sent to the network forensic server and analyzed on this single node [12] only. The *forensic server* is obviously the bottleneck that has to process all the data.

Elimination of all *single points of failure* in proposed, redesigned architecture should remedy this issue. The *actor model* is one of the attractive solutions that solve these problems elegantly and efficiently. The actor model was firstly introduced in 1973 [13]. It comes with a separate unit called *actor*. Actors execute independently and in parallel. They communicate asynchronously via message passing, and their state is otherwise immutable. Actor's behavior determines how to process the incoming message. Actor system is considered to be capable of linear scalability [1].

3. Problem Statement and Solution

Network forensics, is a tedious work that strictly relies on completeness and precision of all undertaken steps to gain a piece of a puzzle that fits together as a shred of evidence. Considering the current speeds of regular users' home network connection(s), an unabridged analysis would require enormous computation resources. Try to imagine, that each network packet would be analysed by many protocol dissectors with a goal to extract for example an acknowledgment of email delivery. To be able to achieve this goal with optimal computation power, I must revisit currently utilized methods and redesign them to work on in distributed environment which brings new challenges to architecture design, algorithms functionality, data synchronization and so on and so forth.

Let's start with an imaginary demonstration. The math is simple, one computer with 1 Gbps NIC (Network Interface Card) that has a relatively simple task to capture traffic during full line load would be required to write to a disk under the constant speed

of 1000Mbps \approx 125 MB/s. Proposed system has to guarantee that no data are lost during the capture. A suspect can simultaneously download and upload data which means that the monitoring device cannot have only one 1 * 1 Gbps NIC, but it needs 2 * 1 Gbps cards, one for uplink, one for downlink. Thus, the required speed of continuous disk writing would be $2 * 125 \text{ MB/s} \approx 250 \text{ MB/s}$. Now, if the requirement is to store the communication for one day, the disk capacity have to be $250 * 60 * 60 * 24 \approx 21.6 \text{ TB}$. This is achievable with commodity hardware, e.g., 2 * 12 TB drives with RAID 0 or 4 * 12 TB with RAID 1+0 with assumed write/read speed of 250 MB/s. But what if only one day is not enough? For typical forensic case, capturing period spawns through weeks or months.

Assuming, one day is what I need. I am tasked to retrieve valuable information from the captured communication. I know that I can perform the write operation on drives, but for analysis purposes, I need to read the data as well. Concurrent read and write operations on the same hard-drives slows down both, starting with data loss in capturing phase, resulting in irretrievable information during the analysis. I cannot afford either one. We can argue, that a "simple" solution would be to double the count of drives to create a performance buffer and let operating system to deal with it. But what if the analytical process requires more computation power at one moment and overwhelm resources in spite of the capturing; it would end up with the data loss. What happens with results of the analysis? It need to store them, and where else than on hard-drives that are under continuous pressure from capturing; it end up with the data loss.

From my previous performance measurements, I know that single computation node is limited and commodity hardware is hardly sufficient to perform all required operations in real-time and over long periods of time. Speed of separation of frames into a conversations which needs a dissection of the network protocols up to the application layer, is roughly 300Mbps [14, pp. 45-51], which is not sufficient. On the other hand, I am confident that application created and optimized for this singular purpose can do the processing faster and breach the 1 Gbps line speed. Nevertheless, I do not believe that single machine solution is capable of doing overall analysis and extraction of information from the application layer. I have to design the solution to be distributed across multiple machines.

The solution is based on the actor model. Each actor represents an independent processing unit. The communication between actors is managed by messaging. The actor has no shared state; thus all actors

138 work in parallel. If actors run on the same node, the
139 message passing has a little overhead compared to a
140 function call or a loop. However, if actors scale over
141 multiple nodes, messages need to be serialized. The
142 serialization process introduces latency and consumes
143 part of processing power.

144 4. Architecture Design

145 Incomplete data provided by unreliable traffic inter-
146 ception can lead to skewed results; some information
147 may be lost, some fabricated by reconstruction process.
148 Keeping these facts in mind, the processing cannot
149 strictly follow RFCs and behave like a *kernel* network
150 stack implementation, but it has to incorporate several
151 heuristics. For example, to fill missing gaps in data,
152 and to consider these fillings during application proto-
153 col processing, or never to join multiple frames into
154 a single conversation unless it passes more advanced
155 heuristics and checks. Network forensic tools which
156 I have worked with do mostly respect RFCs and thus
157 may produce misleading results as already shown [15].

158 I propose a distributed architecture with no single
159 point of failure, composed of commodity hardware
160 that will be capable of linear scalability, and capable
161 of fine resource utilization. See Figure 1 for design
162 details.

163 At the top level, I have divided the entire process
164 into the two main stages:

165 **Data preprocessing** Reconstruction of conversations
166 at the application layer (L7 conversations) from
167 the captured traffic. Each of these conversations
168 holds information about the source and desti-
169 nation endpoints, time stamps and reassembled
170 payloads of exchanged application messages.

171 **Data analysis** Identification of application protocols
172 in reconstructed L7 conversations and subse-
173 quent use of a proper application protocol dis-
174 sector to reconstruct application events from
175 given conversations (e.g., visited web pages,
176 sent emails, ...). The output of this stage is a
177 set of forensic artifacts.

178 *First stage*, data preprocessing, is executed on set
179 of an independent *Reassembler* nodes. Reconstructed
180 L7 conversations from the stream of captured packets
181 can originate from *PCAP files* or can be captured from
182 *the live network interface*.

183 In the most common use case, there is only one
184 source stream (i.e., one *PCAP file*) which I want to
185 analyze. Therefore to utilize all of the *Reassembler*
186 instances, I have to split packets from this stream into a
187 smaller sub-streams, which will be distributed among

online *Reassembler* instances. For this split, I can not
188 use a naive method such as *Round Robin*. *Reassem-*
189 *bler* nodes operate independently of each other and to
190 fully reconstruct L7 conversations (each can consists
191 of multiple packets), they have to obtain all the pieces
192 of the particular L7 conversation. 193

194 Using this naive method, there could occur a situa-
195 tion where half the packets from one L7 conversation
196 will end up in one *Reassembler* node and second half in
197 some other; both nodes would end up with incomplete
198 data, and none of them won't be able to reconstruct
199 the conversation entirely.

200 Solution to this problem is another type of nodes
201 called *L4 Load Balancer*, which will be positioned
202 in front of the *Reassembler* nodes. They will extract
203 source and destination IP addresses and ports from
204 each packet of the source stream and will use them
205 to decide to which instance of *Reassembler* should
206 forward the packet concerning its context. This way,
207 all packets of a particular L7 conversation will always
208 be forwarded to precisely one *Reassembler* instance.
209 The reconstructed L7 conversation will be then stored
210 in a distributed database, ready to be retrieved in the
211 second stage of the execution.

212 In the *second stage*, a subset of reconstructed L7
213 conversations is retrieved from the distributed database
214 (by using manual or automatic selection with specified
215 rules) and delivered to the *Snooper* nodes. They will
216 identify used application protocol and use proper appli-
217 cation protocol dissector module to extract data from
218 the L7 conversation. Extracted data will be stored back
219 into the distributed database.

220 Each instance of a particular node acts as an indi-
221 vidual actor in the system, communicating with other
222 actors by message passing. Thanks to this design, I
223 am able to distribute the computation across multiple
224 machines maintaining the linear scalability.

225 5. Conclusion

226 In this research, I have proposed the system for dis-
227 tributed real-time forensic network traffic analysis up
228 to the application layer capable of processing com-
229 munication at high speed. I intend to create a system
230 based on actor model that scales linearly and is hard-
231 ware independent.

232 My prototype implementation of the proposed sys-
233 tem called NTPAC (Network Traffic Processing &
234 Analysis Cluster) is based on C# actor system library
235 *Akka.NET*. Selection of technologies implementing
236 higher abstractions is essential for fast prototype cre-
237 ation. My preliminary measurements conducted on

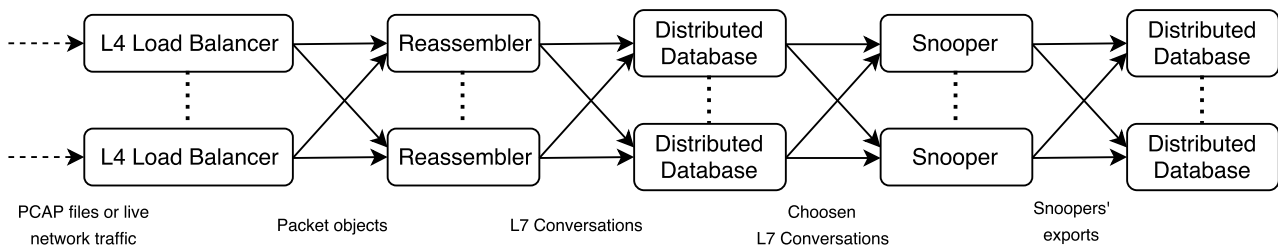


Figure 1. Architecture diagram shows proposed system’s nodes with information flow between them. Interconnections between agents are logical, independent of underlying hardware architecture. The solution is to be deployed on a single node or scale up in a distributed environment.

238 regular workstation¹ show that PCAP files are read
 239 and parsed up to application protocol layer with speed
 240 of 3929Mbps on a single core. In comparison, speed
 241 capture file loading with a single *Capture* actor is
 242 3482Mbps/core. Lastly implemented *L4 Load Bal-*
 243 *ancer* operates at speed of 3447Mbps/2cores. These
 244 are very first measurements without more complex per-
 245 formance optimization on an incomplete system. The
 246 *Reassembler* supporting TCP and UDP with heuris-
 247 tics [15], *Distributed Database* module, supporting
 248 *ArangoDB* or *Cassandra*, are implemented, but not yet
 249 integrated.

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