

# Safeguarding PoS Consensus

## A Comparative Analysis of Proposer Protection Mechanisms Against Various Attack Vectors

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### Proof-of-Stake Consensus

- validators update the chain by **proposing blocks**
- proposers are **randomly selected**
- incentive: rewards, slashable stake

### Ethereum

- proposer selected every 12 seconds (one slot)
- selection made **in advance** to let validators prepare
- results are **publicly known**, even to attackers
- possibly nation-state actors

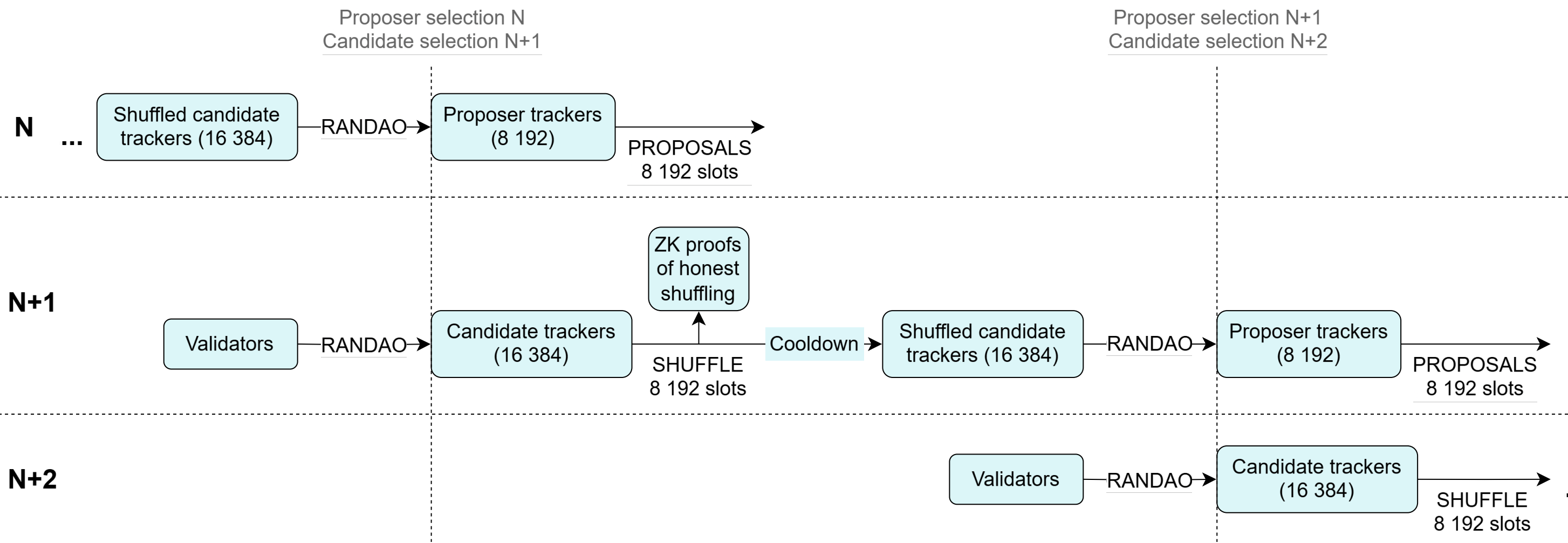
### Simulation

- goal:** measure effectiveness and overhead of protection mechanisms
- protection mechanisms:** Secret Single Leader Election

### Secret Single Leader Election (SSLE)

- Only proposer knows they have been selected
- Whisk:**
  - shuffling
  - hidden identity: BLS G1 point, random secret
  - Zero-Knowledge Proofs
- Homomorphic sortition:**
  - Threshold Fully Homomorphic Encryption
  - hidden identity: encrypted random numbers
  - independent FHE circuits

Fig. 1: The pipeline in the SSLE algorithm Whisk.



### Attack: Malicious DoS

Attacker aims to negatively impact the network by attempting to DoS every proposer.

Fig. 2: Malicious DoS with no protection.



Fig. 3: Malicious DoS with Whisk.

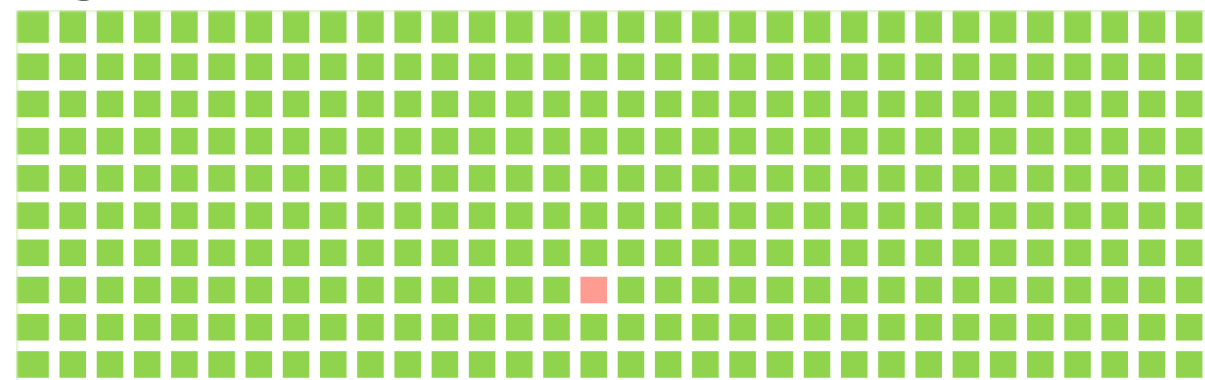
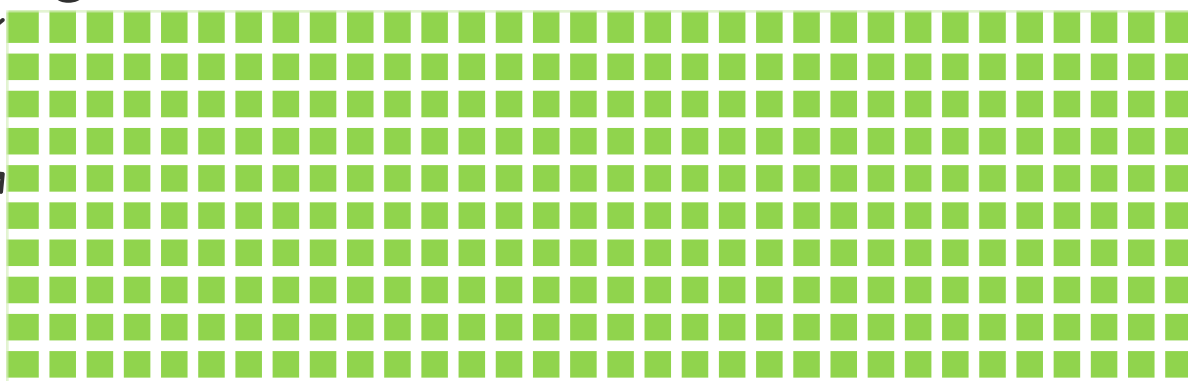


Fig. 4: Advanced malicious DoS with Whisk.



attacker randomly attacks many candidates at once after the shuffling phase (arrows)

only if proposer gets selected twice in a row (very rare - over 1 mil. validators)

### Attack: Censorship

Attacker targets a validator group, preventing them from participating in the consensus.

Fig. 5: Censorship with no protection.

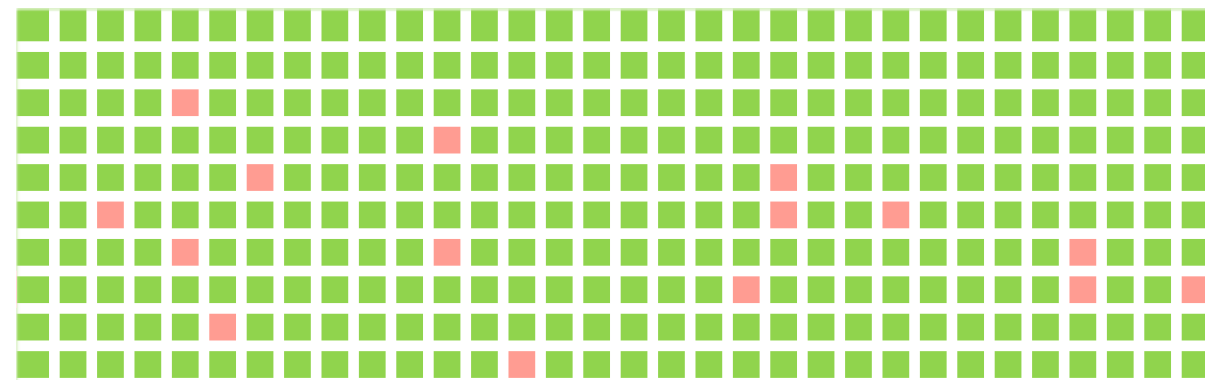


Fig. 6: Censorship with Whisk.

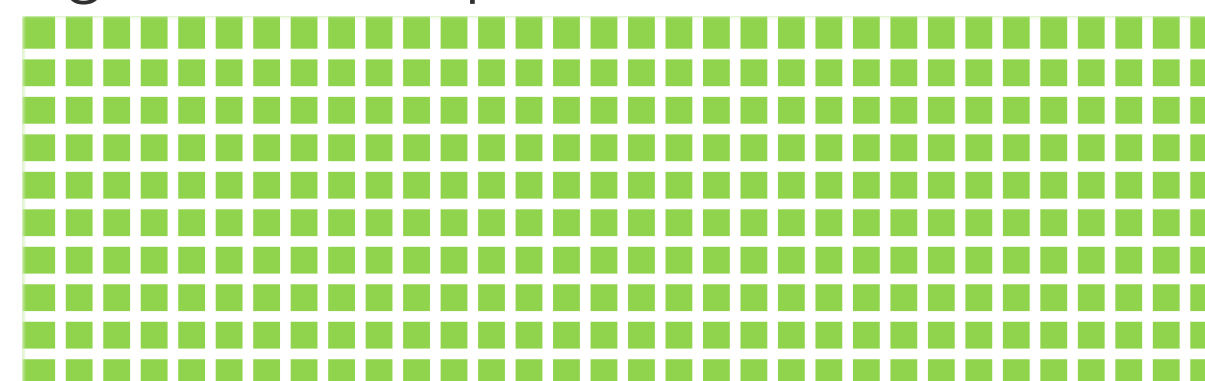
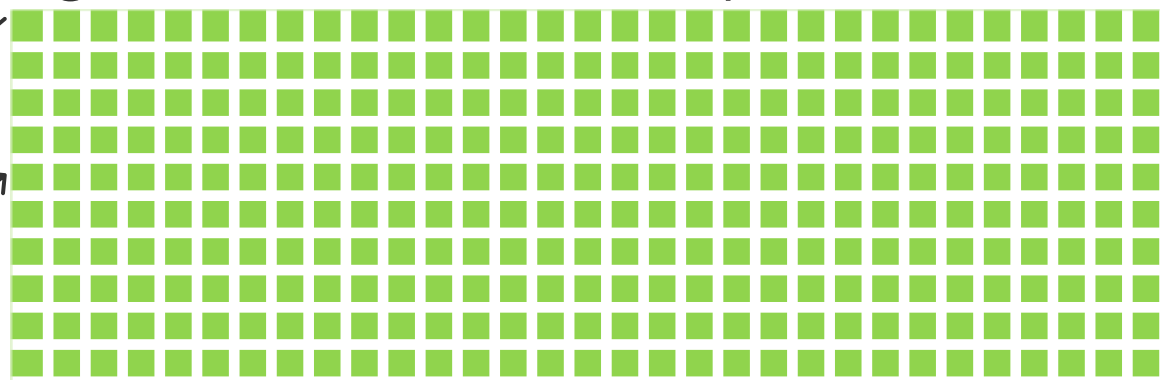


Fig. 7: Advanced censorship with Whisk.



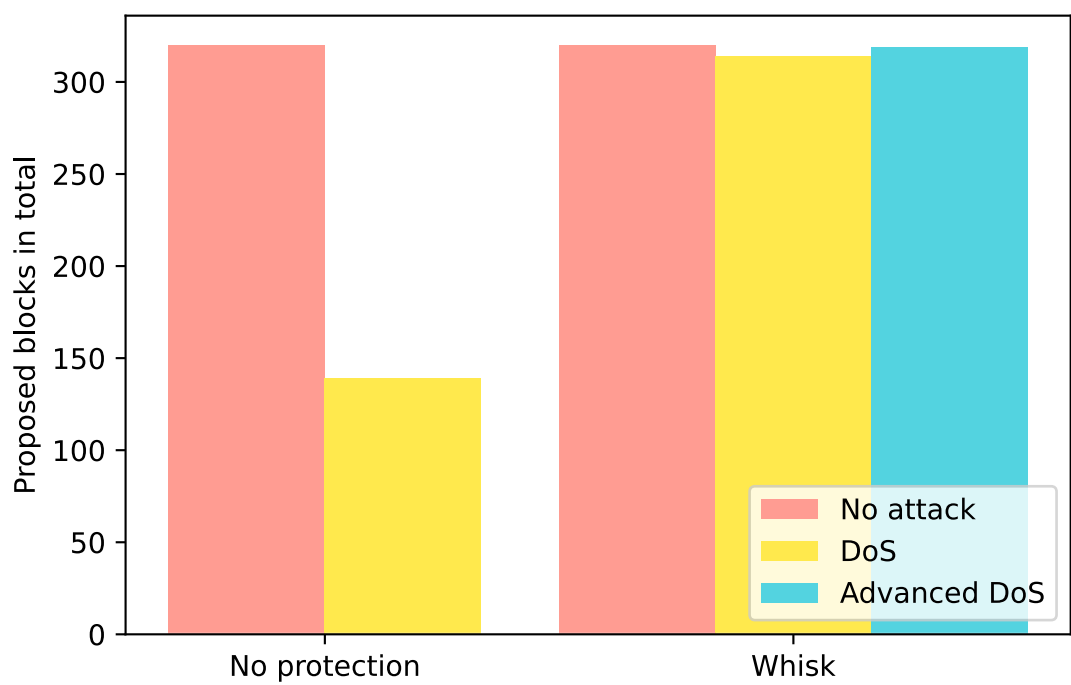
continuous DoS of victims that also were candidates before shuffling, low probability of success

the rarity of DoS success combined with a smaller group of victims

### Protection Results

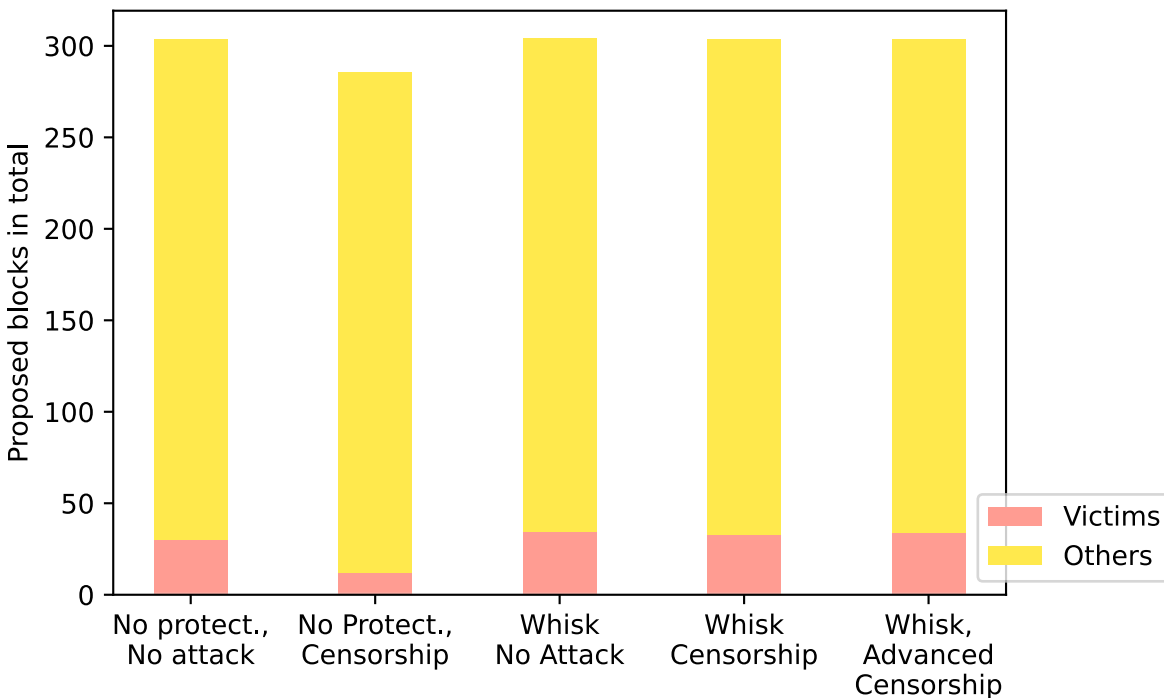
- DoS:
- slots missed: **55.43 % / 1.36 %**
  - proposers affected: **56.63 % / 1.36 %**

Fig. 8: Comparison of proposed blocks during DoS.



- Censorship:
- slots missed: **6.00 % / 0.21 %**
  - victims affected: **61.82 % / 2.19 %**

Fig. 9: Comparison of proposed blocks during censorship.



- Time measurements in one run:
- no protection: **5 160 ms**
  - Whisk: **74 961 ms (14x longer)**